

Problems as Solutions

Unique versus Constructive Existence

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... If a question can be put at all, then it can also be answered.

L. Wittgenstein [*TLP* 6.5]

Overview

Introduction and Folklore
Without Countable Choice
Why the Fan Theorem?
From Brouwer to Gödel

Situation:

(S, d) metric space, usually complete or even compact

$F : S \rightarrow \mathbb{R}_{\geq 0}$ continuous, mostly uniformly continuous

Problem:

If $\inf_{x \in S} F(x) = 0$, is there $\xi \in S$ with $F(\xi) = 0$?

Alternative formulation:

$$\forall \varepsilon > 0 \exists x \in S (F(x) < \varepsilon) \stackrel{?}{\Rightarrow} \exists \xi \in S \forall \varepsilon > 0 (F(\xi) < \varepsilon)$$

approximate solutions $\overset{?}{\rightsquigarrow}$ exact solution

First attempt:

Choose a sequence (x_n) in S with $F(x_n) < 1/n$ for all n .

In general this is an invocation of countable choice.

If S is compact, then (x_n) has a cluster point ξ in S , for which $F(\xi) = 0$.

We thus have used the Bolzano–Weierstraß theorem

BWT *Every sequence in a compact space has a cluster point,*

as one does in the usual proof of the minimum theorem

MIN *Every continuous function on a compact space has a minimum,*

by which we would have solved our problem anyway.

Both BWT and MIN, however, are essentially non–constructive.

There is a construction of $\inf_{x \in S} F(x)$ if F is uniformly continuous, and if S is compact: that is, totally bounded and complete.

So when does our problem still have a constructive solution?

In other words, when does a uniformly continuous function on a compact space have a minimum: that is, attain its infimum?

We say that F has *uniformly at most one* minimum if

$$\forall \delta \exists \varepsilon \forall x, y (F(x) < \varepsilon \wedge F(y) < \varepsilon \Rightarrow d(x, y) < \delta)$$

where δ, ε are positive rationals and x, y points of S .

If F has uniformly at most one minimum, then F has a minimum.

In fact, the sequence (x_n) in S with $F(x_n) < 1/n$ is a Cauchy sequence. Hence if S is complete, then (x_n) has a limit ξ in S , for which $F(\xi) = 0$.

One only needs that S is complete, and F sequentially continuous.

This was used intensively by U. Kohlenbach in the early 1990s. He could even start from a so-called modulus of uniqueness, which gives a modulus of convergence for the sequence (x_n) .

It is a constructive solution, but a priori requires countable choice, by which the given data *is* an element of the completion of S .

We can do the same without choice following F. Richman, who considers the completion of S as the set \hat{S} of all locations on S .

A *location* on S is an $f : S \rightarrow \mathbb{R}$ with $\inf f = 0$ so that

$$|f(x) - f(y)| \leq d(x, y) \leq f(x) + f(y) .$$

The metric on the set \hat{S} of all locations on S is defined by

$$d(f, g) = \sup |f - g| = \inf (f + g) .$$

There is the isometric embedding

$$S \hookrightarrow \hat{S}, z \mapsto d(z, \cdot) ,$$

with which S is dense in \hat{S} , and \hat{S} is complete: that is, $\hat{S} = \hat{\hat{S}}$.

Note that $d(f, z) = f(z)$ for all $f \in \hat{S}$ and $z \in S$.

Every $f \in \hat{S}$ is uniformly continuous, has uniformly at most one minimum, and satisfies

$$f(y) = \lim_{f(x) \rightarrow 0} d(x, y) .$$

The locations on \mathbb{Q} form the Dedekind reals \mathbb{R} , for which $\mathbb{R} = \hat{\mathbb{R}}$.

Each uniformly continuous $\varphi : S \rightarrow T$ extends uniquely to a

$$\hat{\varphi} : \hat{S} \rightarrow \hat{T} \quad \text{with} \quad \hat{\varphi}(f)(y) = \lim_{f(x) \rightarrow 0} d(\varphi(x), y)$$

which is uniformly continuous.

With all this, we can solve our problem without countable choice.

Theorem 1 If $F : S \rightarrow \mathbb{R}$ with $\inf F = 0$ is uniformly continuous and has uniformly at most one minimum, then

$$f(y) = \lim_{F(x) \rightarrow 0} d(x, y)$$

defines an $f \in \hat{S}$ with $\hat{F}(f) = 0$ for $\hat{F} : \hat{S} \rightarrow \mathbb{R}$.

$$\left[d(\hat{F}(f), 0) \stackrel{\check{=}}{=} \hat{F}(f)(0) \stackrel{\check{=}}{=} \lim_{f(x) \rightarrow 0} \overbrace{d(F(x), 0)}^{F(x)} \stackrel{(!)}{=} 0 \right]$$

If $S = \mathbb{R}$ and $F(x) = |x - \xi|^k$ with $k \geq 1$ and $\xi \in \mathbb{R}$, then

$$f(y) = \lim_{x \rightarrow \xi} d(x, y) = d(\xi, y) = |y - \xi|.$$

If F is already a location, then $f = F$.

Brouwer's fan theorem can equivalently be put as

FAN *Every decidable binary tree without infinite path is finite.*

This is the contrapositive of weak König's lemma

WKL *Every infinite decidable binary tree has an infinite path,*

which is essentially non-constructive.

The logical form of FAN and WKL is

$$\text{FAN} \quad \forall \alpha \exists n P(\bar{\alpha}n) \Rightarrow \exists n \forall \alpha P(\bar{\alpha}n)$$

$$\text{WKL} \quad \forall n \exists \alpha P(\bar{\alpha}n) \Rightarrow \exists \alpha \forall n P(\bar{\alpha}n)$$

where $n \in \mathbb{N}$ and P is a decidable property of the finite initial segments

$$\bar{\alpha}n = (\alpha(0), \dots, \alpha(n-1))$$

of infinite binary sequences α : that is, the points of the Cantor space.

We say that F as before has *at most one* minimum if

$$\forall x, y (d(x, y) > 0 \Rightarrow F(x) > 0 \vee F(y) > 0) .$$

With J. Berger and D. Bridges we found out in 2003 that

MIN! *Every uniformly continuous function on a compact space that has at most one minimum possesses a minimum*

is equivalent to FAN, but it was unclear why FAN has to do with this.

In 2005 J. Berger and H. Ishihara carried this over to other cases: e.g., with Brouwer's fixed point theorem in place of the minimum theorem.

All these equivalents of FAN have the form

If a problem on a compact space has approximate solutions and at most one solution, then it has an exact solution.

It was still unclear why FAN occurred in this context.

The positivity principle

POS *Every uniformly continuous function on a compact space that attains only positive values has a positive infimum*

is the equivalent of FAN that was needed to prove MIN!.

With FAN in the form of POS we now know why it occurred there.

Theorem 2 FAN is equivalent to

UAM *If a uniformly continuous function on a compact space has at most one minimum, then it has uniformly at most one minimum.*

This sharpens FAN \Rightarrow MIN!, because UAM \Rightarrow MIN!.

The implications MIN! \Rightarrow FAN and FAN \Rightarrow POS have been proved before by J. Berger, D. Bridges, H. Ishihara, and the author.

To show the missing link POS \Rightarrow UAM, we assume that S is compact, and $F : S \rightarrow \mathbb{R}_{\geq 0}$ uniformly continuous.

Note first that $d(x, y) \geq \delta$ defines a compact subset of $S \times S$ for all but countably many values of δ .

The condition “ F has at most one minimum” can be put as

$$\forall x, y (\exists \delta d(x, y) \geq \delta \Rightarrow F(x) + F(y) > 0) ,$$

which is equivalent to

$$\forall \delta \forall x, y (d(x, y) \geq \delta \Rightarrow F(x) + F(y) > 0) .$$

By POS, this implies

$$\forall \delta \exists \varepsilon \forall x, y (d(x, y) \geq \delta \Rightarrow F(x) + F(y) \geq \varepsilon) ,$$

which is equivalent to “ F has uniformly at most one minimum”.

A *modulus of uniqueness* is a function $\mu : \delta \mapsto \varepsilon$ with

$$\forall \delta \forall x, y (F(x) < \mu(\delta) \wedge F(y) < \mu(\delta) \Rightarrow d(x, y) < \delta) .$$

Such a μ can be obtained from the corresponding $\forall \delta \exists \varepsilon$ -statement “ F has uniformly at most one minimum” by countable choice.

Corollary *In the presence of countable choice, FAN is equivalent to*

MUN *If a uniformly continuous function on a compact space has at most one minimum, then it has a modulus of uniqueness.*

For S bounded, U. Kohlenbach obtained a modulus of uniqueness from

$$(*) \quad \forall x, y (F(x) = 0 \wedge F(y) = 0 \Rightarrow d(x, y) = 0)$$

by his monotone functional interpretation (m.f.i.) going back to Gödel.

Condition $(*)$ is the contrapositive of “ F has at most one minimum”, and has the same m.f.i. whenever S is bounded.

In the presence of countable choice, FAN is thus equivalent to an instance of the m.f.i.: namely, to MUN.