

Algebraic Topology

(Castro-Urdiales tutorial)

I. Combinatorial Topology

```
;; Clock
Computing
<TnPr <TnPr
End of computing.

;; Clock -> 2002-01-17, 19h 25m 36s.
Computing the boundary of the generator 19 (dimension 7) :
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>> <<Abar>> <<Abar>>
End of computing.
```

Homology in dimension 6 :

Component Z/12Z

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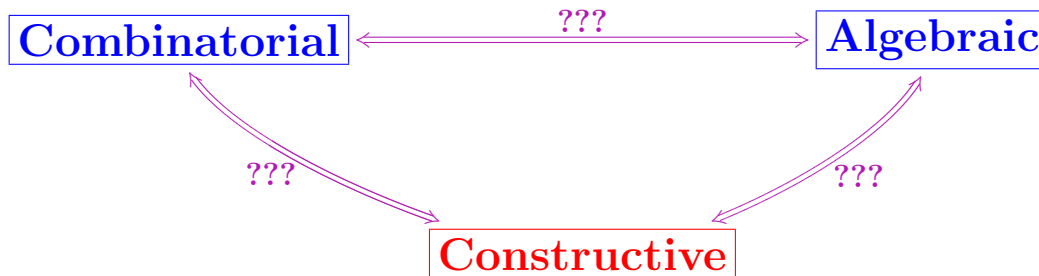
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;; Clock -> 2002-01-17, 19h 27m 15s
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Castro-Urdiales, January 9-13, 2006*

Tutorial plan:

1. Combinatorial Topology.
2. Homological Algebra.
3. **Constructive** Algebraic Topology.
4. **Implemented** Algebraic Topology.

Essential terminological problem:



“General” topological spaces

cannot be directly installed in a computer.

A combinatorial translation is necessary.

Main methods:

1. Simplicial complexes.
2. Simplicial sets.

Warning: Simplicial sets much more complex (!)

but much more powerful than simplicial complexes.

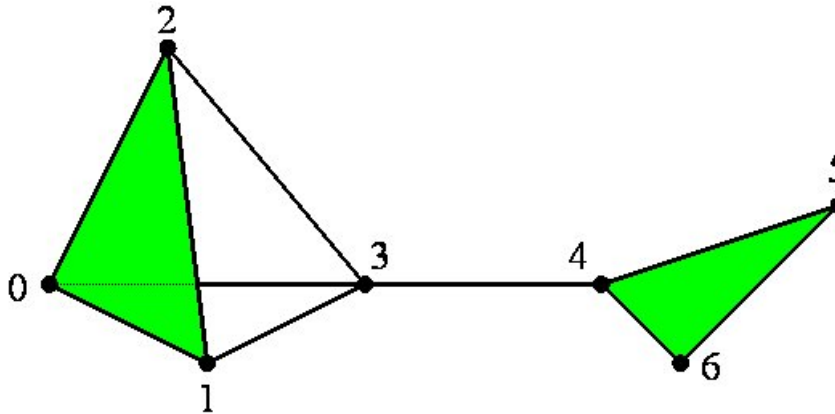
Simplicial complex $K = (V, S)$ where:

1. $V = \text{set} = \text{set of vertices of } K$;
2. $S \in \mathcal{P}(\mathcal{P}_*^f(V))$ (= set of simplices) satisfying:
 - (a) $\sigma \in S \Rightarrow \sigma = \text{non-empty finite set of vertices}$;
 - (b) $\{v\} \in S$ for all $v \in V$;
 - (c) $\{(\sigma \in S) \text{ and } (\emptyset \neq \sigma' \subset \sigma)\} \Rightarrow (\sigma' \in S)$.

Notes:

1. V may be **infinite** ($\Rightarrow S$ **infinite**).
2. $\forall \sigma \in S, \sigma$ is finite.

Example:



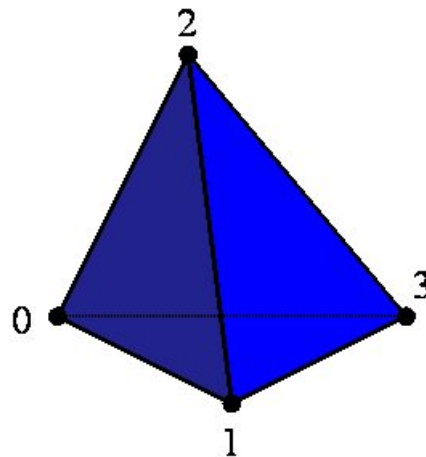
$$V = \{0, 1, 2, 3, 4, 5, 6\}$$

$$S = \left\{ \begin{array}{l} \{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \\ \{0, 1\}, \{0, 2\}, \{0, 3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{3, 4\}, \\ \{4, 5\}, \{4, 6\}, \{5, 6\}, \{0, 1, 2\}, \{4, 5, 6\} \end{array} \right\}$$

Drawbacks of simplicial complexes.

Example: 2-sphere :

Needs 4 vertices, 6 edges, 4 triangles.



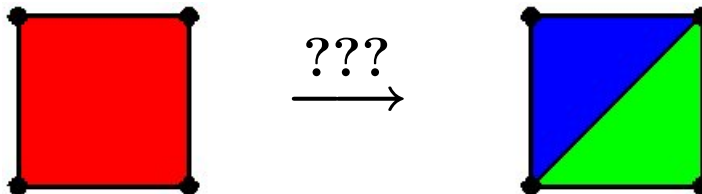
The simplicial set model needs only

1 vertex + 1 “triangle”

but an infinite number of degenerate simplices...

Product?

$$\Delta^1 \times \Delta^1 = I \times I?$$



In general, constructions are difficult

with simplicial complexes.

Main support for the notion of **simplicial set**:

The Δ category.

Objects: $\underline{0} = \{0\}, \underline{1} = \{0, 1\}, \dots, \underline{m} = \{0, 1, \dots, m\}, \dots$

Morphisms:

$$\Delta(\underline{m}, \underline{n}) = \{\alpha : \underline{m} \rightarrow \underline{n} \text{ st } (k \leq l \Rightarrow \alpha(k) \leq \alpha(l))\}.$$

Examples:

$$\partial_{\boxed{i}}^m = \begin{array}{c} 0 \longrightarrow 0 \\ 1 \longrightarrow 1 \\ \vdots \text{ : } \vdots \\ i-1 \longrightarrow i-1 \\ i \longrightarrow \boxed{i} \\ \vdots \text{ : } \vdots \\ m-2 \longrightarrow m-2 \\ m-1 \longrightarrow m-1 \\ m \longrightarrow m \end{array}$$

$$\eta_{\boxed{i}}^m = \begin{array}{c} 0 \longrightarrow 0 \\ 1 \longrightarrow 1 \\ \vdots \text{ : } \vdots \\ i-1 \longrightarrow i-1 \\ i \longrightarrow \boxed{i} \\ i+1 \longrightarrow \vdots \\ \vdots \text{ : } \vdots \\ m \longrightarrow m-1 \\ m+1 \longrightarrow m \end{array}$$

Proposition: Any Δ -morphism $\alpha : \underline{m} \nearrow \underline{n}$

has a **unique** expression:

$$\alpha = \partial_{i_1}^n \cdots \partial_{i_{n-p}}^{p+1} \eta_{j_{m-p}}^p \cdots \eta_{j_1}^{m-1}$$

with $i_1 > \cdots > i_{n-p}$ and $j_{m-p} < \cdots < j_1$ if

$$\alpha : \underline{m} \nearrow \underline{p} \nearrow \underline{n}$$

is the unique Δ -factorisation

through a **surjection** and an **injection**.

Example:

$$\Delta(\underline{4}, \underline{5}) \ni \alpha = \begin{array}{|l} 0 \rightarrow 0 \\ 1 \rightarrow 0 \\ 2 \rightarrow 0 \\ 3 \rightarrow 2 \\ 4 \rightarrow 2 \end{array} = \partial_5^5 \partial_4^4 \partial_3^3 \partial_1^2 \eta_0^1 \eta_1^2 \eta_3^3$$

Definition: A **simplicial set** is

a contravariant functor $X : \Delta \rightarrow \mathbf{Sets}$.

Example: K simplicial complex $\mapsto X_K$ simplicial set.

$K = (V, S)$ with V totally ordered.

$\mapsto X_K = (\{X_m\}_{m \in \text{Ob}(\Delta)}, \{X_\alpha\}_{\alpha \in \text{Mfp}(\Delta)})$ with:

$X_m = \{(v_0 \leq \dots \leq v_m) \text{ st } \{v_0, \dots, v_m\} \in S\};$

$= \{\chi : \underline{m} \nearrow V \text{ monotone and } S\text{-compatible}\};$

$\Delta(\underline{n}, \underline{m}) \ni [\alpha : \underline{n} \nearrow \underline{m}] \mapsto [X_\alpha : X_m \rightarrow X_n : \chi \mapsto \chi \circ \alpha].$

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Definition : Let $X = (\{X_m\}, \{X_\alpha\})$ be a **simplicial set**.

The **geometric realization** of X is defined by:

$$|X| = \left(\coprod_m X_m \times \Delta^m \right) / \sim$$

with $(x, \alpha_* t) \sim (\alpha^* x, t)$ for $x \in X_m$, $t \in \Delta^n$, $\alpha \in \Delta(\underline{n}, \underline{m})$.

Proposition: Every **point** of $|X|$ has

a **unique representant** $(x, t) \in X_m \times \Delta_m$ satisfying:

1. x is a **non-degenerate simplex**;
2. $t \in \text{Int } \Delta^m$.

What is a **non-degenerate simplex**?

Definition: A simplex $\sigma \in X_m$ is **degenerate**
if $\sigma = \alpha^* \sigma'$ where $\sigma' \in X_n$ with $n < m$.

Degenerate = Necessary **consequence**
of another simplex in lower dimension.

Proposition: Every simplex σ has a **unique expression**
 $\sigma = \alpha^* \sigma'$ with σ' **non-degenerate** and α surjective.

Every **degenerate simplex** is a **degeneracy**
of a **unique non-degenerate simplex**.

Corollary: $|X| = \left(\coprod_m X_m^{ND} \times \Delta^m \right) / \sim_{ND}$.

\Rightarrow In fact, only the **non-degenerate simplices**
are **visible** in the **geometric realization**.

Example 1: Standard **3-simplex** D (solid tetrahedron).

$$m \in \mathbb{N} \Rightarrow D_m = \Delta(\underline{m}, \underline{3})$$

$$\begin{aligned} [\alpha \in \Delta(\underline{n}, \underline{m})] + [\sigma \in D_m = \Delta(\underline{m}, \underline{3})] \\ \Rightarrow [D_\alpha(\sigma) = \alpha^*(\sigma) = \sigma \circ \alpha \in D_n = \Delta(\underline{n}, \underline{3})]. \end{aligned}$$

$$\Rightarrow [\sigma \in D_m = \Delta(\underline{m}, \underline{3}) \text{ non-degenerate}] \Leftrightarrow [\sigma \text{ injective}].$$

$$\Rightarrow \left\{ \begin{array}{l} \#D_0^{ND} = 4 \\ \#D_1^{ND} = 6 \\ \#D_2^{ND} = 4 \\ \#D_3^{ND} = 1 \\ \#D_m^{ND} = 0 \text{ if } m \geq 4 \end{array} \right. \Rightarrow \text{standard model of } \Delta_3.$$

Example 2: $X_m = \{*_m\} \coprod \Delta^{\text{srj}}(\underline{m}, \underline{2})$.

Given $\sigma \in X_m$ and $\alpha \in \Delta(\underline{n}, \underline{m})$,

$$X_\alpha(\sigma) = \alpha^*(\sigma) = ?$$

1. $[\sigma = *_m] \Rightarrow [\alpha^*(\sigma) = *_n]$;
2. $\sigma \in \Delta^{\text{srj}}(\underline{m}, \underline{2})$, consider $\sigma \circ \alpha$:
 - (a) If $\sigma \circ \alpha \in \Delta^{\text{srj}}(\underline{n}, \underline{2})$, then $\alpha^*(\sigma) = \sigma \circ \alpha$;
 - (b) Otherwise $\alpha^*(\sigma) = *_n$.

Exercise: $X_0^{ND} = \{*_0\}$, $X_2^{ND} = \{\text{id}_2\}$,

others X_m^{ND} are empty.

Corollary: $|X| = S^2$.

Example 3: $G = \{0, 1\} = \mathbb{Z}/2\mathbb{Z}$.

$$\begin{cases} X_m = G^m = \{\underline{m} \rightarrow G\}. \\ [\sigma : \underline{m} \rightarrow G] + [\alpha : \underline{n} \rightarrow \underline{m}] \Rightarrow [\alpha^*(\sigma) = \sigma \circ \alpha : \underline{n} \rightarrow G]. \end{cases}$$

$$|X| = ???$$

Exercise: $\sigma : \underline{m} \rightarrow G$ is **degenerate**

if and only if two **consecutive images** are **equal**.

Corollary: $X_m^{ND} = \{(0, 1, 0, 1, \dots), (1, 0, 1, 0, \dots)\}$.

$$\#X_0^{ND} = 2 \Rightarrow |X|_0 = 2 \text{ points} = S^0.$$

$$\#X_1^{ND} = 2 \Rightarrow |X|_1 = 2 \text{ edges attached to 2 points} = S^1.$$

$$\#X_2^{ND} = 2 \Rightarrow |X|_2 = 2 \text{ "triangles" attached to } S^1 = S^2.$$

... ..

$$\Rightarrow |X| = \varinjlim S^n = S^\infty.$$

Example 4: Let us consider again the previous example X .

A **canonical group action** $\mathbb{Z}_2 \times X \rightarrow X$ is defined.

Dividing X by this action produces a new simplicial set Y .

In particular $\#X_m^{ND} = 2 \Rightarrow \#Y_m^{ND} = 1$.

And $|Y| = S^\infty / \mathbb{Z}_2 = P^\infty \mathbb{R} =$ **infinite real projective space**.

Example 5: Replacing \mathbb{Z}_2 by any **discrete group** G produces
in the same way the **fundamental space** $K(G, 1)$.

Definition: $K(G, 1)$ is the **base space** of a **Galoisian covering**
where the **structural group** is G
and the **total space** is **topologically contractible**.

Product construction for simplicial sets.

$X = (\{X_m\}, \{X_\alpha\})$, $Y = (\{Y_m\}, \{Y_\alpha\})$ two simplicial sets.

$Z = X \times Y = ???$

Simple and natural definition:

$Z = X \times Y$ defined by $Z = (\{Z_m\}, \{Z_\alpha\})$ with:

$$Z_m = X_m \times Y_m$$

If $\Delta(\underline{n}, \underline{m}) \ni \alpha : \underline{n} \nearrow \underline{m}$:

$$Z_\alpha: X_m \times Y_m \xrightarrow{X_\alpha \times Y_\alpha} X_n \times Y_n$$

Example: $I \times I = \Delta^1 \times \Delta^1 = ???$

$$Z_m = \Delta(\underline{m}, \underline{1}) \times \Delta(\underline{m}, \underline{1})$$

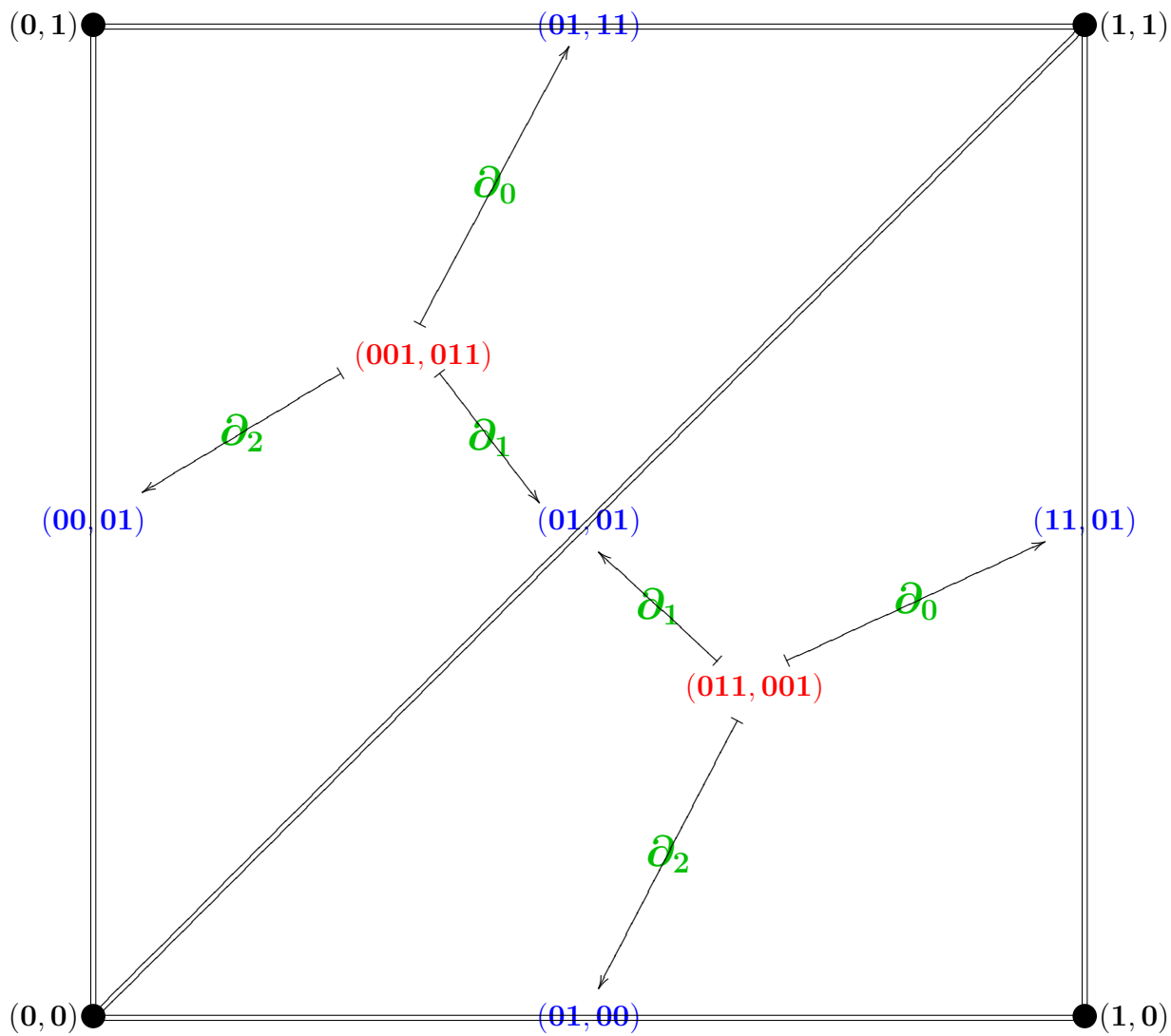
$$(011, 001) \in Z_2$$

$$\partial_0 : 01 \nearrow 012 \quad \partial_0(011, 001) = (11, 01)$$

$$\partial_1 : 01 \nearrow 012 \quad \partial_1(011, 001) = (01, 01)$$

$$\partial_2 : 01 \nearrow 012 \quad \partial_2(011, 001) = (01, 00)$$

Picture ???



Twisted Products.

$B = (\{B_m\}, \{B_\alpha\})$ with:

$$\begin{cases} B_0^{ND} = \{*\} \\ B_1^{ND} = \{s_1\} \\ B_m^{ND} = \emptyset \text{ if } m \geq 2 \end{cases} \Rightarrow B = \text{standard model of } S^1.$$

$G = (\{G_m\}, \{G_\alpha\})$

with $G_m = \mathbb{Z}$ ($\forall m \in \mathbb{N}$) and $G_\alpha = \text{id}_{\mathbb{Z}}$ ($\forall \alpha \in \Delta(\underline{m}, \underline{n})$).

Exercise: $|G| = \mathbb{Z}$.

$\Rightarrow B \times G = S^1 \times \mathbb{Z} = \text{stack of circles.}$

In particular: $\partial_0(s_1, k_1) = (*, k_0)$.

Definition: $B = \text{simplicial set}$ + $G = \text{simplicial group}$.

A twisting function $\tau : B \rightarrow G$

is a family $(\tau_m : B_m \rightarrow G_{m-1})_{m \geq 1}$ satisfying:

$\partial_i \tau(b) = \tau \partial_{i+1}(b) \quad (i \geq 1)$	$\partial_0 \tau(b) = (\tau \partial_0(b))^{-1} \tau \partial_1(b)$
$\eta_i \tau(b) = \tau \eta_{i+1}(b) \quad (i \geq 0)$	$\tau \eta_0(b) = e_{m-1} \quad (b \in B_m)$

New face operators for $B \times_{\tau} G$:

$\partial_i(b, g) = (\partial_i b, \partial_i g) \quad (i \geq 1)$	$\partial_0(b, g) = (\partial_0 b, \tau(b) \cdot \partial_0(b))$
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Example 1: $B = S^1$, $G = \mathbb{Z}$, $\tau_1(s_1) = \mathbf{0}_0 \in G_0$
 $\Rightarrow B \times_{\tau} G = S^1 \times \mathbb{Z} = \mathbf{trivial}$ product.

In particular $\partial_0(s_1, \tau(s_1).k_1) = (*, k_0)$.

Example 2: $B = S^1$, $G = \mathbb{Z}$, $\tau'_1(s_1) = \mathbf{1}_0 \in G_0$
 $\Rightarrow B \times_{\tau} G = S^1 \times \mathbb{Z} = \mathbf{twisted}$ product.

Now $\partial_0(s_1, k_1) = (*, \tau'(s_1).k_0) = (*, (k + 1)_0)$.

Daniel Kan's fantastic work (\sim 1960 – 1980).

Every “standard” natural topological **construction process**
has a **translation** in the **simplicial world**.

Frequently the **translation** is even “**better**”.

Typical example. The **loop space construction** in ordinary topology gives only an **H -space** (= **group up to homotopy**).

Kan's loop space construction produces a **genuine** **simplicial group**, playing an **essential role** in Algebraic Topology.

Conclusion: **Simplicial world** = **Paradise!** ???

Simplicial group model for the unit circle $S^1 \subset \mathbb{C}^2$?

$$S^1 = X = (\{X_m\}, \{X_\alpha\}) = ???$$

Eilenberg-MacLane solution (1955):

$$X_m = \mathbb{Z}^m + \text{appropriate } X_\alpha\text{'s}$$

$\Rightarrow X$ highly infinite in every dimension.

Kan (1960) \Rightarrow Eilenberg-MacLane model is the minimal one!

\Rightarrow Simplicial world = Paradise or Hell???

The END

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Homology in dimension 6 :

Component Z/12Z

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